7. Separating Hyperplane Theorems II

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Mathematics II

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Farkas' Lemma

Proposition 7.16 (Farkas' Lemma) Let $A \in \mathbb{R}^{M \times N}$ and $b \in \mathbb{R}^{N}$. The following conditions are equivalent: 1. There exists $x \in \mathbb{R}^{M}$ such that $A^{T}x = b$ and $x \ge 0$.

2. For any
$$y \in \mathbb{R}^N$$
, if $Ay \ge 0$, then $b^{\mathrm{T}}y \ge 0$.

For proof, we will use the following:

Lemma 7.17 $\{A^{\mathrm{T}}x \in \mathbb{R}^N \mid x \in \mathbb{R}^M_+\}$ is a closed set.

Proof of Farkas' Lemma

▶ (1) \Rightarrow (2): Immediate.

(2)
$$\Rightarrow$$
 (1):

Suppose that (1) does not hold.

Let $K = \{A^{\mathrm{T}}x \in \mathbb{R}^N \mid x \in \mathbb{R}^M_+\}$. Then $b \notin K$.

▶ *K* is convex, and by Lemma 7.17 is closed.

▶ Then by the Separating Hyperplane Theorem, there exist $y \in \mathbb{R}^N$ with $y \neq 0$ and $c \in \mathbb{R}$ such that

 $y^{\mathrm{T}}b < c \leq y^{\mathrm{T}}z$ for all $z \in K$,

and therefore, $y^{\mathrm{T}}b < \inf_{z \in K} y^{\mathrm{T}}z$.

- Since K is a cone, it follows that $\inf_{z \in K} y^{\mathrm{T}} z = 0$. (\rightarrow Homework)
- ▶ Thus we have $y^{\mathrm{T}}b < 0$, and $y^{\mathrm{T}}A^{\mathrm{T}}x \ge 0$ for all $x \ge 0$, which implies that $y^{\mathrm{T}}A^{\mathrm{T}} \ge 0^{\mathrm{T}}$.

Proof of Lemma 7.17

Show that $K = \{A^{\mathrm{T}}x \in \mathbb{R}^N \mid x \in \mathbb{R}^M_+\}$ is closed.

- Denote the column vectors in A^{T} by a^1, \ldots, a^M , so that $K = \operatorname{Cone}\{a^1, \ldots, a^M\}.$
- ▶ Let $\{z^m\}$ be a sequence in K, and suppose that $z^m \to \overline{z}$. We want to show that $\overline{z} \in K$.
- By Carathéodory's Theorem, for each m, z^m is written as a conic combination of a linearly independent subset of {a¹,...,a^M}.
- ► Since there are finitely many such subsets, there is a linearly independent subset {aⁱ¹,...,a^{iL}} such that infinitely many elements of {z^m} are written as its conic combinations.
- ▶ Denote $B = (a^{i_1} \cdots a^{i_L}) \in \mathbb{R}^{N \times L}$, and denote the corresponding subsequence again by $\{z^m\}$.

- Denote $z^m = B\lambda^m$, where $\lambda^m \in \mathbb{R}^L_+$.
- We have $B^{\mathrm{T}}z^m = B^{\mathrm{T}}B\lambda^m$, where $B^{\mathrm{T}}B \in \mathbb{R}^{L \times L}$ is non-singular:

• Let $B^{\mathrm{T}}Bx = 0$.

- Then $x^{\mathrm{T}}B^{\mathrm{T}}Bx = 0$, where $x^{\mathrm{T}}B^{\mathrm{T}}Bx = ||Bx||^2$.
- Therefore, $x^{\mathrm{T}}B^{\mathrm{T}}Bx = 0$ if and only if Bx = 0.
- Since the columns of B are linearly independent, this holds if and only if x = 0.
- Therefore, we have $\lambda^m = (B^T B)^{-1} B^T z^m$.
- ▶ By the continuity of $(B^{\mathrm{T}}B)^{-1}B^{\mathrm{T}}z$ in z, λ^m converges to $\bar{\lambda} = (B^{\mathrm{T}}B)^{-1}B^{\mathrm{T}}\bar{z}$, where $\bar{\lambda} \in \mathbb{R}^L_+$.
- Thus, by the continuity of $B\lambda$ in λ , we have $\bar{z} = \lim_{m \to \infty} B\lambda^m = B\bar{\lambda}$, so that $\bar{z} \in K$.

Proposition 7.18 (Farkas' Lemma: Inequality version) Let $A \in \mathbb{R}^{M \times N}$ and $b \in \mathbb{R}^N$. The following conditions are equivalent:

- 1. There exists $x \in \mathbb{R}^M$ such that $A^{\mathrm{T}}x \leq b$ and $x \geq 0$.
- 2. For any $y \in \mathbb{R}^N$, if $y \ge 0$ and $Ay \ge 0$, then $b^T y \ge 0$.

Condition (1) is equivalent to:

There exist $x \in \mathbb{R}^M$ and $z \in \mathbb{R}^N$ such that $x \ge 0$, $z \ge 0$, and $A^{\mathrm{T}}x + z = b$, or $\begin{pmatrix} A^{\mathrm{T}} & I \end{pmatrix} \begin{pmatrix} x \\ z \end{pmatrix} = b$.

By Farkas' Lemma, this is equivalent to:

For any
$$y\in \mathbb{R}^N$$
, if $inom{A}{I}y\geq 0$, then $b^{\mathrm{T}}y\geq 0$,

or, if $y \ge 0$ and $Ay \ge 0$, then $b^{\mathrm{T}}y \ge 0$ (condition (2)).

Linear Programming Let $A \in \mathbb{R}^{K \times N}$, $f \in \mathbb{R}^N$, $c \in \mathbb{R}^K$.

Primal problem:

$$\begin{array}{ll} (\mathsf{P}) & \max_{x \in \mathbb{R}^N} & f^{\mathsf{T}}x \\ & \mathsf{s.t.} & Ax \leq c \\ & x \geq 0. \end{array}$$

Dual problem:

$$\begin{array}{ll} \text{(D)} & \min_{\lambda \in \mathbb{R}^{K}} \ c^{\mathrm{T}}\lambda \\ & \text{s. t.} \quad A^{\mathrm{T}}\lambda \geq f \\ & \lambda \geq 0. \end{array}$$

The Lagrangians for the two problems coincide (the nonnegativity constraints aside):

$$L(x,\lambda) = f^{\mathrm{T}}x - \lambda^{\mathrm{T}}(Ax - c) = c^{\mathrm{T}}\lambda - x^{\mathrm{T}}(A^{\mathrm{T}}\lambda - f).$$

Weak Duality

Proposition 7.19 If $x \in \mathbb{R}^N$ and $\lambda \in \mathbb{R}^K$ are feasible for (P) and (D), respectively, then $f^{\mathrm{T}}x \leq c^{\mathrm{T}}\lambda$.

Proof

▶ If
$$x \in \mathbb{R}^N$$
 and $\lambda \in \mathbb{R}^K$ are feasible for (P) and (D), then
 $f^{\mathrm{T}}x \leq (A^{\mathrm{T}}\lambda)^{\mathrm{T}}x = \lambda^{\mathrm{T}}(Ax) \leq \lambda^{\mathrm{T}}c.$

Therefore, if $\bar{x} \in \mathbb{R}^N$ and $\bar{\lambda} \in \mathbb{R}^K$ are feasible and if $f^T \bar{x} = c^T \bar{\lambda}$, then \bar{x} and $\bar{\lambda}$ are solutions to (P) and (D), respectively.

Strong Duality

Proposition 7.20

Suppose that both (P) and (D) are feasible. Then both (P) and (D) have solutions, and

 $\max\{f^{\mathrm{T}}x\mid Ax\leq c,\ x\geq 0\}=\min\{c^{\mathrm{T}}\lambda\mid A^{\mathrm{T}}\lambda\geq f,\ \lambda\geq 0\}.$

Suppose that (P) and (D) are feasible.

We want to show that there exist $x \in \mathbb{R}^N$ and $\lambda \in \mathbb{R}^K$ such that $Ax \leq c$, $A^{\mathrm{T}}\lambda \geq f$, $f^{\mathrm{T}}x \geq c^{\mathrm{T}}\lambda$, $x \geq 0$, and $\lambda \geq 0$, or

$$\begin{pmatrix} A & O \\ O & -A^{\mathrm{T}} \\ -f^{\mathrm{T}} & c^{\mathrm{T}} \end{pmatrix} \begin{pmatrix} x \\ \lambda \end{pmatrix} \leq \begin{pmatrix} c \\ -f \\ 0 \end{pmatrix}, \ x \ge 0, \ \lambda \ge 0.$$

▶ By Farkas' Lemma (inequality version; Proposition 7.18), this is equivalent to the condition that for all $p \in \mathbb{R}^K$, $q \in \mathbb{R}^N$, and $r \in \mathbb{R}$,

$$\begin{pmatrix} p^{\mathrm{T}} & q^{\mathrm{T}} & r \end{pmatrix} \begin{pmatrix} A & O \\ O & -A^{\mathrm{T}} \\ -f^{\mathrm{T}} & c^{\mathrm{T}} \end{pmatrix} \ge 0, \ p \ge 0, \ q \ge 0, \ r \ge 0$$
$$\Rightarrow \begin{pmatrix} p^{\mathrm{T}} & q^{\mathrm{T}} & r \end{pmatrix} \begin{pmatrix} c \\ -f \\ 0 \end{pmatrix} \ge 0.$$



(1)
$$A^{\mathrm{T}}p \ge rf, \ Aq \le rc, \ p \ge 0, \ q \ge 0, \ r \ge 0$$

implies

(2)
$$c^{\mathrm{T}}p - f^{\mathrm{T}}q \ge 0.$$

We want to show that this holds whenever (P) and (D) are feasible.

- For r > 0, (1) implies that q/r and p/r are feasible solutions to (P) and (D), so that we have
 c^Tp − f^Tq = r[c^T(p/r) − f^T(q/r)] ≥ 0 by Weak Duality.
- For r = 0, let x and λ be feasible solutions to (P) and (D). From (1), we have

$$c^{\mathrm{T}}p - f^{\mathrm{T}}q \ge x^{\mathrm{T}}A^{\mathrm{T}}p - \lambda^{\mathrm{T}}Aq \ge 0.$$

Strong Duality

Proposition 7.21

1. Suppose that (D) has a solution.

Then (P) has a solution, and

 $\max\{f^{\mathrm{T}}x \mid Ax \le c, \ x \ge 0\} = \min\{c^{\mathrm{T}}\lambda \mid A^{\mathrm{T}}\lambda \ge f, \ \lambda \ge 0\}.$

2. Suppose that (P) has a solution. Then (D) has a solution, and

 $\max\{f^{\mathrm{T}}x\mid Ax\leq c,\;x\geq 0\}=\min\{c^{\mathrm{T}}\lambda\mid A^{\mathrm{T}}\lambda\geq f,\;\lambda\geq 0\}.$

Suppose that (D) has a solution.

In light of Proposition 7.20, it suffices to show that (P) has a feasible solution.

• Let $\lambda^* \in \mathbb{R}^K$ be a solution to (D).

To apply Farkas' Lemma (Proposition 7.18), let $z \in \mathbb{R}^K$ be such that $A^{\mathrm{T}}z \geq 0$ and $z \geq 0$.

▶ Then $\lambda^* + z \ge 0$, and $A^T(\lambda^* + z) = A^T\lambda^* + A^Tz \ge f$, which means that $\lambda^* + z$ is feasible in (D).

• Therefore, by the optimality of λ^* , we have $0 \le c^{\mathrm{T}}(\lambda^* + z) - c^{\mathrm{T}}\lambda^* = c^{\mathrm{T}}z$.

▶ By Proposition 7.18, there exists $x \in \mathbb{R}^N$ such that $Ax \leq c$ and $x \geq 0$.

Variants of Farkas' Lemma

Proposition 7.22 (Gale's Theorem) Let $A \in \mathbb{R}^{M \times N}$ and $b \in \mathbb{R}^N$. The following conditions are equivalent:

1. There exists $x \in \mathbb{R}^M$ such that $A^{\mathrm{T}}x \leq b$.

2. For any $y \in \mathbb{R}^N$, if $y \ge 0$ and Ay = 0, then $b^T y \ge 0$.

Condition (1) is equivalent to:

There exist $z^1 \in \mathbb{R}^M$ and $z^2 \in \mathbb{R}^M$ such that $z^1 \ge 0$, $z^2 \ge 0$, and $A^{\mathrm{T}}(z^1 - z^2) \le b$, or $\begin{pmatrix} A^{\mathrm{T}} & -A^{\mathrm{T}} \end{pmatrix} \begin{pmatrix} z^1 \\ z^2 \end{pmatrix} \le b$.

 By Farkas' Lemma (inequality version; Proposition 7.18), this is equivalent to:

For any $y \in \mathbb{R}^N$, if $y \ge 0$ and $\begin{pmatrix} A \\ -A \end{pmatrix} y \ge 0$, then $b^T y \ge 0$, or, if $y \ge 0$ and Ay = 0, then $b^T y \ge 0$ (condition (2)).

Proposition 7.23 (Gordan's Theorem) Let $A \in \mathbb{R}^{M \times N}$. The following conditions are equivalent:

- 1. There exists $x \in \mathbb{R}^M$ such that $A^{\mathrm{T}}x \gg 0$.
- 2. For any $y \in \mathbb{R}^N$, if $y \ge 0$ and Ay = 0, then y = 0.

Condition (1) is equivalent to:

There exists $x \in \mathbb{R}^M$ such that $-A^{\mathrm{T}}x \leq -\mathbf{1}$.

 By Gale's Theorem (Proposition 7.22), this is equivalent to: For any y ∈ ℝ^N, if y ≥ 0 and (-A)y = 0, then (-1^T)y ≥ 0, or y ≥ 0 and Ay = 0, then y = 0 (condition (2)). Proposition 7.24 (Ville/von Neumann-Morgenstern I) Let $A \in \mathbb{R}^{M \times N}$. The following conditions are equivalent:

- 1. There exists $x \in \mathbb{R}^M$ such that $A^T x \gg 0$ and $x \gg 0$.
- 2. For any $y \in \mathbb{R}^N$, if $y \ge 0$ and $Ay \le 0$, then y = 0.

Variants of Farkas' Lemma

▶ In fact, "there exists $x \in \mathbb{R}^M$ such that $A^T x \gg 0$ and $x \gg 0$ " is equivalent to "there exists $x \in \mathbb{R}^M$ such that $A^T x \gg 0$ and $x \ge 0$ ".

• Given an $x \ge 0$ in the latter, consider $x + \varepsilon \mathbf{1}$ for sufficiently small $\varepsilon > 0$.

Proposition 7.25 (Ville/von Neumann-Morgenstern II) Let $A \in \mathbb{R}^{M \times N}$. The following conditions are equivalent:

1. There exists $x \in \mathbb{R}^M$ such that $A^T x \gg 0$ and $x \ge 0$.

2. For any $y \in \mathbb{R}^N$, if $y \ge 0$ and $Ay \le 0$, then y = 0.

Proof of Proposition 7.24

Condition (1) is equivalent to:

There exists
$$x \in \mathbb{R}^M$$
 such that $\begin{pmatrix} A^{\mathrm{T}} \\ I \end{pmatrix} x \gg 0.$

▶ By Gordan's Theorem (Proposition 7.23), this is equivalent to: For any $y \in \mathbb{R}^N$ and $z \in \mathbb{R}^M$, if $y \ge 0$, $z \ge 0$, and $\begin{pmatrix} A & I \end{pmatrix} \begin{pmatrix} y \\ z \end{pmatrix} = 0$, then $\begin{pmatrix} y \\ z \end{pmatrix} = 0$.

This is equivalent to:

For any $y \in \mathbb{R}^N$, if $y \ge 0$ and $Ay \le 0$, then y = 0 (condition (2)).

Variants of Farkas' Lemma

Proposition 7.26 Let $A \in \mathbb{R}^{M \times N}$. The following conditions are equivalent:

- 1. There exists $x \in \mathbb{R}^M$ such that $A^T x \leq 0$ and $x \gg 0$.
- 2. For any $y \in \mathbb{R}^N$, if $y \ge 0$ and $Ay \ge 0$, then Ay = 0.

Condition (1) is equivalent to:

There exists
$$x \in \mathbb{R}^M$$
 such that $\begin{pmatrix} A^{\mathrm{T}} \\ -I \end{pmatrix} x \leq \begin{pmatrix} 0 \\ -\mathbf{1} \end{pmatrix}$.

▶ By Gale's Theorem (Proposition 7.22), this is equivalent to:
For any
$$y \in \mathbb{R}^N$$
 and $z \in \mathbb{R}^M$,
if $y \ge 0$, $z \ge 0$, and $\begin{pmatrix} A & -I \end{pmatrix} \begin{pmatrix} y \\ z \end{pmatrix} = 0$, then
 $\begin{pmatrix} 0 & -\mathbf{1}^T \end{pmatrix} \begin{pmatrix} y \\ z \end{pmatrix} \ge 0.$

This is equivalent to:

For any $y \in \mathbb{R}^N$, if $y \ge 0$ and $Ay \ge 0$, then Ay = 0 (condition (2)). Proposition 7.27 (Stiemke's Lemma) Let $A \in \mathbb{R}^{M \times N}$. The following conditions are equivalent:

1. There exists $x \in \mathbb{R}^M$ such that $A^{\mathrm{T}}x = 0$ and $x \gg 0$.

2. For any $y \in \mathbb{R}^N$, if $Ay \ge 0$, then Ay = 0.

Condition (1) is equivalent to:

There exists
$$x \in \mathbb{R}^M$$
 such that $x \gg 0$ and $\begin{pmatrix} A^T \\ -A^T \end{pmatrix} x \leq 0$.

▶ By Proposition 7.26, this is equivalent to:
For any
$$y \in \mathbb{R}^N$$
 and $z \in \mathbb{R}^M$,
if $y \ge 0$, $z \ge 0$, and $\begin{pmatrix} A & -A \end{pmatrix} \begin{pmatrix} y \\ z \end{pmatrix} \ge 0$, then
 $\begin{pmatrix} A & -A \end{pmatrix} \begin{pmatrix} y \\ z \end{pmatrix} = 0.$

This is equivalent to:

For any $y \in \mathbb{R}^N$, if $Ay \ge 0$, then Ay = 0 (condition (2)).

Variants of Farkas' Lemma

Proposition 7.28 (Motzkin's Theorem) Let $B \in \mathbb{R}^{M \times N}$, $C \in \mathbb{R}^{M \times K}$, $D \in \mathbb{R}^{M \times L}$. The following conditions are equivalent:

- 1. There exists no $x \in \mathbb{R}^M$ such that $B^T x \gg 0$, $C^T x \ge 0$, and $D^T x = 0$.
- 2. There exist $y_1 \in \mathbb{R}^N$, $y_2 \in \mathbb{R}^K$, and $y_3 \in \mathbb{R}^L$ such that $By_1 + Cy_2 + Dy_3 = 0$, $y_1 \ge 0$, $y_1 \ne 0$, and $y_2 \ge 0$.

- Proved using Farkas' Lemma.
- Proposition 7.23 (Gordan's Theorem), Propositions 7.24-7.25 (Ville's Theorem), Proposition 7.26, and Proposition 7.27 (Stiemke's Lemma) are all special cases of this theorem.

Efficient Production under Linear Technology

For the production set $Y \subset \mathbb{R}^N$, $y \in Y$ is *efficient* if there is no $y' \in Y$ such that $y' \ge y$ and $y' \ne y$.

Proposition 7.29 Let $Y = \{y \in \mathbb{R}^N \mid Ay \le b\}$ for some $A \in \mathbb{R}^{M \times N}$ and $b \in \mathbb{R}^M$. Then $\overline{y} \in Y$ is efficient if and only if there exists $p \gg 0$ such that

 $p \cdot \bar{y} \ge p \cdot y$ for all $y \in Y$.

▶ The "if" part holds for general set *Y*.

► The "if" part:

If \bar{y} is not efficient, i.e., $y' - \bar{y} \ge 0$, $\neq 0$ for some $y' \in Y$, then for any $p \gg 0$, we have $(y' - \bar{y})p > 0$ or $y'p > y^*p$.

The "only if" part:

Suppose that $\bar{y} \in Y$ is efficient.

▶ Write
$$A = \begin{pmatrix} A^1 \\ A^2 \end{pmatrix}$$
 and $b = \begin{pmatrix} b^1 \\ b^2 \end{pmatrix}$ such that
 $A^1 \bar{y} = b^1, \qquad A^2 \bar{y} \ll b^2,$
where $A^k \in \mathbb{R}^{M_k \times N}$, $b^k \in \mathbb{R}^{M_k}$, $k = 1, 2$, and $M_1 + M_2 = M$.
▶ By the efficiency of \bar{y} , $M_1 \ge 1$.

▶ By the efficiency of \bar{y} , there exists no $z \in \mathbb{R}^N$ such that $A^1z \leq 0, z \geq 0, z \neq 0.$

If there exists such z, then $A(\bar{y} + \varepsilon z) \leq b$ for sufficiently small $\varepsilon > 0$, where $\bar{y} + \varepsilon z \geqq \bar{y}$.

By Proposition 7.25 (Ville's Theorem), there exists x ∈ ℝ^{M1} such that (A¹)^Tx ≫ 0 and x ≥ 0.
 Let p = (A¹)^Tx (≫ 0).

▶ Then for any $y \in Y$ (where $A^1y \leq b^1$), we have

$$p \cdot \bar{y} = x \cdot A^1 \bar{y} = x \cdot b^1,$$

$$p \cdot y = x \cdot A^1 y \le x \cdot b^1,$$

as desired.

Strict Dominance and Never Best Response

Consider a two-player normal form game:

- S₁ = {1,..., M}: set of pure strategies of player 1 (M ≥ 2)
 S₂ = {1,..., N}: set of pure strategies of player 2 (N ≥ 2)
- $\Delta(S_1) = \{x \in \mathbb{R}^M_+ \mid x_1 + \ldots + x_M = 1\}$: set of mixed strategies of player 1 $\Delta(S_2) = \{y \in \mathbb{R}^N_+ \mid y_1 + \ldots + y_N = 1\}$: set of mixed strategies of player 2
- From player 1's point of view, Δ(S₂) is interpreted as the set of 1's *beliefs* over 2's strategies.
- Pure strategy $m \in S_1$ is identified with $e_m \in \Delta(S_1)$, the *m*th unit vector of \mathbb{R}^M .

Payoff matrix for player 1:

$$U = \begin{pmatrix} u_{11} & \cdots & u_{1N} \\ \vdots & \ddots & \vdots \\ u_{M1} & \cdots & u_{MN} \end{pmatrix} \in \mathbb{R}^{M \times N}$$

(We only consider the incentives of player 1.)

- $e_m^{\mathrm{T}} Uy \cdots$ payoff from $m \in S_1$ against $y \in \Delta(S_2)$
- $x^{\mathrm{T}}Uy \cdots$ payoff from $x \in \Delta(S_1)$ against $y \in \Delta(S_2)$

- $m \in S_1$ is a best response to $y \in \Delta(S_2)$ if $e_m^{\mathrm{T}} U y \ge e_{\ell}^{\mathrm{T}} U y$ for all $\ell \in S_1$.
- *m* ∈ S₁ is a *never best response* if it is not a best response to any *y* ∈ Δ(S₂).
- $m \in S_1$ is strictly dominated if there exists $x \in \Delta(S_1)$ such that $e_m^{\mathrm{T}} U e_n < x^{\mathrm{T}} U e_n$ for all $n \in S_2$.

Proposition 7.30

In a two-player normal form game, $m \in S_1$ is a never best response if and only if it is strictly dominated.

The result extends straightforwardly to (finite) games with more than two players if best response is defined with respect to *correlated* beliefs over opponents' strategies.

Let

$$\tilde{U} = \begin{pmatrix} u_{11} - u_{m1} & \cdots & u_{1N} - u_{mN} \\ \vdots & \ddots & \vdots \\ u_{M1} - u_{m1} & \cdots & u_{MN} - u_{mN} \end{pmatrix}.$$

▶
$$m \in S_1$$
 is a never best response
 \iff there exists no $y \ge 0$, $y \ne 0$, such that $\tilde{U}y \le 0$
 \iff if $y \ge 0$ and $\tilde{U}y \le 0$, then $y = 0$.

• $m \in S_1$ is strictly dominated \iff there exists $x \ge 0$, $x \ne 0$, such that $x^{\mathrm{T}} \tilde{U} \gg 0$.

▶ By Ville's Theorem (Proposition 7.25), these are equivalent.

Weak Dominance and Never Best Response

• $m \in S_1$ is weakly dominated if there exists $x \in \Delta(S_1)$ such that

•
$$e_m^{\mathrm{T}} U e_n \leq x^{\mathrm{T}} U e_n$$
 for all $n \in S_2$, and

•
$$e_m^{\mathrm{T}}Ue_n < x^{\mathrm{T}}Ue_n$$
 for some $n \in S_2$.

Proposition 7.31

In a two-player normal form game, $m \in S_1$ is a best response to some totally mixed strategy $y \in \Delta(S_2)$ if and only if it is not weakly dominated.

Again let

$$\tilde{U} = \begin{pmatrix} u_{11} - u_{m1} & \cdots & u_{1N} - u_{mN} \\ \vdots & \ddots & \vdots \\ u_{M1} - u_{m1} & \cdots & u_{MN} - u_{mN} \end{pmatrix}$$

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• $m \in S_1$ is a best response to some totally mixed strategy \iff there exists $y \gg 0$ such that $\tilde{U}y \leq 0$.

$$\begin{array}{l} \blacktriangleright \ m \in S_1 \text{ is not weakly dominated} \\ \iff \ \text{there exists no } x \geq 0, \ x \neq 0, \ \text{such that } x^{\mathrm{T}} \tilde{U} \geqq 0 \\ \iff \ \text{if } x \geq 0 \ \text{and } x^{\mathrm{T}} \tilde{U} \geq 0, \ \text{then } x^{\mathrm{T}} \tilde{U} = 0. \end{array}$$

By Proposition 7.26, these are equivalent.